

Pattern Synonyms

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Example

```
-- Datatype definition
data Type = TyApp String [ Type]

-- Functions abstract construction
tyInt :: Type
tyInt = TyApp "Int" []

mkFunTy :: Type -> Type -> Type
mkFunTy t u = TyApp "->" [t, u]

plusTy :: Type
plusTy = tyInt 'mkFunTy' tyInt 'mkFunTy' tyInt
```

Example

```
-- Datatype definition
data Type = TyApp String [ Type]

-- Pattern synonym abstracts matching
pattern FunTy :: Type → Type → Type
pattern FunTy t u = TyApp "->" [t, u]
funArgTys :: Type → ([ Type], Type)
funArgTys (FunTy t u) = case funArgTys u of
  (ts, r) → (t : ts, r)
funArgTys t = ([], t)
```

Pattern synonyms

- ▶ Goal: bring function-like abstraction to pattern matching
- ▶ Touches all parts of the GHC frontend:
 - ▶ Parser, renamer
 - ▶ Typechecker
 - ▶ Desugarer
 - ▶ Interface files
- ▶ No backend changes needed
- ▶ Our paper shows the breadth; here we show some depth:
 - ▶ Typechecking
 - ▶ Desugaring

Pattern types

The Typing Principle

It should be possible to determine whether a use of P is well-typed based only on P 's type, without reference to P 's definition.

What is the type of a pattern?

Pattern types

Scrutinee type:

pattern $P1 = True$

pattern $P2 = Just True$

pattern $P3 = Nothing$

Parametric patterns:

pattern $P4\ x = Just\ x$

Pattern types

Scrutinee type:

```
pattern P1 :: Bool  
pattern P1 = True  
pattern P2 :: Maybe Bool  
pattern P2 = Just True  
pattern P3 ::  $\forall a.$ Maybe a  
pattern P3 = Nothing
```

Parametric patterns:

```
pattern P4 x = Just x
```


Pattern types

Scrutinee type:

```
pattern P1 :: Bool  
pattern P1 = True  
pattern P2 :: Maybe Bool  
pattern P2 = Just True  
pattern P3 ::  $\forall a.$ Maybe a  
pattern P3 = Nothing
```

Parametric patterns:

```
pattern P4 ::  $\forall a.$ a  $\rightarrow$  Maybe a  
pattern P4 x = Just x
```

Pattern types: required constraints

Required constraints:

pattern $P5 = 42$

pattern $P6 = (show \rightarrow "foo")$

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pattern $P5 :: \forall a. (Num\ a, Eq\ a) \Rightarrow a$

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Pattern types: required constraints

Required constraints:

pattern $P5 :: \forall a. (Num\ a, Eq\ a) \Rightarrow a$

pattern $P5 = 42$

pattern $P6 :: \forall a. (Show\ a) \Rightarrow a$

pattern $P6 = (show \rightarrow "foo")$

Pattern types: existentials and provided constraints

data T *a* **where**

$MkT :: (Eq\ a, Show\ b) \Rightarrow a \rightarrow a \rightarrow b \rightarrow T\ a$

$f\ (MkT\ x\ y\ v) = \mathbf{if}\ x \equiv y\ \mathbf{then}\ Just\ (show\ v)\ \mathbf{else}\ Nothing$

Matching on MkT brings in scope

- the type b
- $(Eq\ a, Show\ b)$

allowing (\equiv) and $show$ to be used on the right-hand side

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$f :: T\ a \rightarrow Maybe\ String$

Pattern types: existentials and provided constraints

data T a **where**

$MkT :: (Eq\ a, Show\ b) \Rightarrow a \rightarrow a \rightarrow b \rightarrow T\ a$

pattern $P\ x\ y\ v = MkT\ x\ y\ v$

Matching on P brings in scope

- the type b
- $(Eq\ a, Show\ b)$

Pattern types: existentials and provided constraints

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pattern $P\ x\ y\ v = MkT\ x\ y\ v$

Matching on P brings in scope

- the type b
- $(Eq\ a, Show\ b)$

pattern $P :: \forall a. () \Rightarrow \forall b. (Eq\ a, Show\ b) \Rightarrow a \rightarrow a \rightarrow b \rightarrow T\ a$

Pattern types: required & provided constraints

data T a **where**

$MkT :: (Eq\ a, Show\ b) \Rightarrow a \rightarrow b \rightarrow T\ a$

pattern $P\ v = MkT\ 1\ v$

- ▶ Matching on P **requires** $(Eq\ a, Num\ a)$
- ▶ Matching on P **provides** $(Eq\ a, Show\ b)$

Pattern types: required & provided constraints

data $T\ a$ **where**

$MkT :: (Eq\ a, Show\ b) \Rightarrow a \rightarrow b \rightarrow T\ a$

pattern $P\ v = MkT\ 1\ v$

- ▶ Matching on P **requires** $(Eq\ a, Num\ a)$
- ▶ Matching on P **provides** $(Eq\ a, Show\ b)$

pattern $P :: \forall a. (Num\ a) \Rightarrow \forall b. (Eq\ a, Show\ b) \Rightarrow b \rightarrow T\ a$

Pattern synonym types

Pattern synonym types are fully specified on six axes:

1. Universally bound type variables *univ*
2. The required context *Req* (with *univ* in scope)
3. The scrutinee type *t* (with *univ* in scope)

Surface syntax for pattern synonym type signatures:

```
pattern P ::  $\forall univ. Req \Rightarrow t$ 
```

Pattern synonym types

Pattern synonym types are fully specified on six axes:

1. Universally bound type variables *univ*
2. The required context *Req* (with *univ* in scope)
3. The scrutinee type *t* (with *univ* in scope)
4. Existentially bound type variables *ex*
5. The provided context *Prov* (with *univ* and *ex* in scope)
6. The types of parameters *t1*, *t2*, ... (with *univ* and *ex* in scope)

Surface syntax for pattern synonym type signatures:

```
pattern P ::  $\forall univ. Req \Rightarrow \forall ex. Prov \Rightarrow t1 \rightarrow t2 \rightarrow \dots \rightarrow t$ 
```

Desugaring

Desugaring

pattern $P :: \forall a. (\text{Num } a) \Rightarrow \forall b. (\text{Eq } a, \text{Show } b) \Rightarrow$
 $b \rightarrow T a$
pattern $P \ v = \text{MkT } 3 \ v$

$\$mP :: \forall r a. (\text{Num } a) \Rightarrow$
 $T a \rightarrow$
 $(\forall b. (\text{Eq } a, \text{Show } b) \Rightarrow b \rightarrow r) \rightarrow r \rightarrow r$
 $\$mP \ x \ sk \ fk = \text{case } x \ \text{of}$
 $\text{MkT } 3 \ v \rightarrow sk \ v$
 $_ \rightarrow fk$

Can be represented in existing GHC Core

- ▶ No changes needed anywhere downstream
- ▶ Exported, linked, and potentially inlined just like any other function
- ▶ Synthetic *Void#* parameter can be used to prevent incorrect strictness when *r* is unboxed

Desugaring

Semantics is different from macro-substitution!

The full (potentially nested) structure of the pattern synonym is matched first:

```
f1 :: [Bool] → Bool
f1 [True] = True
f1 _      = False
```

```
f1 [⊥, ⊥] = ⊥
f1 (False : ⊥) = False
```

```
pattern Single a = [a]
f2 :: [Bool] → Bool
f2 (Single True) = True
f2 _              = False
```

```
f2 [⊥, ⊥] = False
f2 (False : ⊥) = ⊥
```


Conclusion

- ▶ Proposal added in 2011
- ▶ New in GHC 7.8 in 2014
- ▶ Incremental improvements in GHC 7.10 and 8.0
- ▶ Used in 72 packages on Hackage (as of June 2016)

Check our paper for more. . .

- ▶ Lots of examples
- ▶ Pattern synonym directionality
- ▶ Record syntax support
- ▶ Importing/exporting
- ▶ Shorthand syntax for pattern synonym signatures
- ▶ Formalisation of pattern types (in the extended version)

<http://unsafePerform.IO/patsyn>